

Syllable Structure and Model Theory (Strother-Garcia, 2019)

- rule-based account and OT models produce structures that are unattested in natural languages
→ we need another, more restrictive formalism

1 Word Model and Model Theories

- used for representing words (strings), where each letter is associated with a particular position
- **model theories** define classes of word models that share a common **signature**
- model theory \mathfrak{M} has the signature $\langle \mathfrak{D}; \mathfrak{R}; \mathfrak{F} \rangle$, where:

– \mathfrak{D} is the set of positions:

- * unordered
- * functions and relations determine the positions
- * usually represented with a set of natural numbers
- * example: *ball*

$$\mathfrak{D} \stackrel{\text{def}}{=} \{1, 2, 3, 4\}$$

– \mathfrak{R} is the set that includes a relation R_σ for every symbol σ in the alphabet Σ .

- * evaluated as Boolean expressions
- * example: *ball*

$$\begin{aligned} \Sigma &= \{b, a, l\} \\ \mathfrak{R} &\stackrel{\text{def}}{=} \{R_b, R_a, R_l\}, \text{ where } R_i \text{ are unary relations} \\ 2 \in R_a \text{ or } R_a(2) \text{ or } a(2) &\rightarrow \text{position 2 is labeled as } a, \text{ etc.} \end{aligned}$$

– \mathfrak{F} is the set of unary functions

- * map domain positions to domain positions
- * example: *ball*

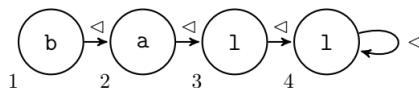
$$\begin{aligned} \mathfrak{F} &\stackrel{\text{def}}{=} \{pred(x), succ(x)\} \\ succ(1) &= 2 \text{ or } 1 \triangleleft 2 \end{aligned}$$

- * **total** function yields output for every position of the domain; because $succ(x) = x + 1$, the final position in the string will also be its own successor ($succ(n) = n$)

1.1 The Successor Model Theory

$$\mathfrak{M}^\triangleleft \stackrel{\text{def}}{=} \langle \mathbb{N}; \{R_\sigma | \sigma \in \Sigma\}; \{pred(x), succ(x)\} \rangle$$

Visual representation of $\mathcal{M}_{ball}^\triangleleft$:



Word model:

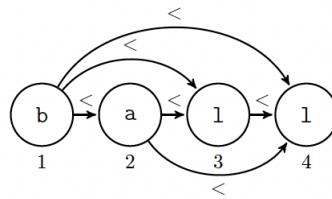
$$\begin{aligned} \mathcal{M}_{ball}^{\triangleleft} &\stackrel{\text{def}}{=} \langle \mathcal{D}; \{R_a, R_b, R_l\}; \{pred(x), succ(x)\} \rangle \\ \mathcal{D} &\stackrel{\text{def}}{=} \{1, 2, 3, 4\} \\ R_a &\stackrel{\text{def}}{=} \{2\} \\ R_b &\stackrel{\text{def}}{=} \{1\} \\ R_l &\stackrel{\text{def}}{=} \{3, 4\} \\ succ(x) &\stackrel{\text{def}}{=} \begin{cases} 2 \Leftrightarrow x = 1 \\ 3 \Leftrightarrow x = 2 \\ 4 \Leftrightarrow x \in \{3, 4\} \end{cases} \\ pred(x) &\stackrel{\text{def}}{=} \begin{cases} 1 \Leftrightarrow x \in \{1, 2\} \\ 2 \Leftrightarrow x = 3 \\ 3 \Leftrightarrow x = 4 \end{cases} \end{aligned}$$

1.2 The Precedence Model Theory

- the **domain** is a set of natural numbers
- unary relations and **general precedence relation** – two positions (binary relation) are ordered with respect to one another with a possibility of intervening positions, e.g. $R_{<}(x, y)$, $3 < 5$
- because of one-to-many relation, precedence cannot be encoded as a **function** $\rightarrow \mathfrak{F} \stackrel{\text{def}}{=} \emptyset$

$$\mathfrak{M}^< \stackrel{\text{def}}{=} \langle \mathbb{N}; \{R_{<}, R_{\sigma} | \sigma \in \Sigma\}; \emptyset \rangle$$

Visual representation of $\mathcal{M}_{ball}^<$:



Word model:

$$\begin{aligned} \mathcal{M}_{ball}^< &\stackrel{\text{def}}{=} \langle \mathcal{D}; \{R_{<}, R_a, R_b, R_l\}; \emptyset \rangle \\ \mathcal{D} &\stackrel{\text{def}}{=} \{1, 2, 3, 4\} \\ R_{<} &\stackrel{\text{def}}{=} \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 4 \rangle\} \\ R_a &\stackrel{\text{def}}{=} \{2\} \\ R_b &\stackrel{\text{def}}{=} \{1\} \\ R_l &\stackrel{\text{def}}{=} \{3, 4\} \end{aligned}$$

2 Enriching Conventional Word Models

- **conventional** models are structured in a very simple way with an ordered sequence of alphabetic characters
- for linguistic purposes we would want to allow for more complex representation of a segment, e.g. phonological features

2.1 Enriching the Alphabet

- in a conventional model, sets of labeling relations for any two symbols are **disjoint**: each position belongs to a unary relation

$$\boxed{(\forall x, y \in \Sigma)[R_x \cap R_y = \emptyset]}$$

- we enrich the model by allowing more than one label per position
- it allows us to maintain similarities between segments, e.g. we want a model to be able to represent similarity between [b] and [p] such that both segments are composed of [labial, stop] features bundle
- let \mathcal{F} be a set of primitive features; and then $\Sigma = \mathcal{F}$

$$\boxed{\mathcal{F} \stackrel{\text{def}}{=} \{voice, cons, high, lab, alv, post, pal, vel, uv, phar, glot, stop, fric, nas, approx, lat\}}$$

- for each feature f in \mathcal{F} , there is a unary relation $R_f \in \mathfrak{R}$. Let \mathcal{R}_f be the set of such relations

$$\boxed{\mathcal{R}_f \stackrel{\text{def}}{=} \{R_f | f \in \mathcal{F}\}}$$

– example: *ball*

$$\boxed{\begin{array}{l} R_{voice}(x) \stackrel{\text{def}}{=} \{1\} \\ R_{labial}(x) \stackrel{\text{def}}{=} \{1\} \\ R_{stop}(x) \stackrel{\text{def}}{=} \{1\} \end{array}}$$

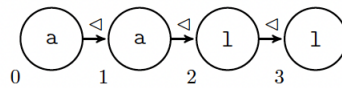
2.2 Enriching the Structure

- conventional models allow for only linear order whether it is general or immediate
- some linguistic representations, such as tone or syllable, require hierarchical representation which relies on **dominance** (which Jeff will talk about)

3 Graph Transductions

- in order to represent input-output mapping, we can use graphs transductions
- in order to map word model \mathfrak{M}^A (input) to another word model \mathfrak{M}^B (output) we must define a set of formulas, one for each relation R and function F in \mathfrak{M}^B

Visual representation of $\Gamma_{b \rightarrow a}(\mathcal{M}_{ball}^{\triangleleft})$:



Definition of the transduction $\Gamma_{b \rightarrow a}(\mathcal{M}_{ball}^{\triangleleft})$, where:

- w indicates an output position
- in the example below both input and output are share the same model theory, however it is not a requirement for such transductions

$$R_a(x^\omega) \stackrel{\text{def}}{=} R_a(x) \vee R_b(x)$$

$$R_b(x^\omega) \stackrel{\text{def}}{=} \text{FALSE}$$

$$R_l(x^\omega) \stackrel{\text{def}}{=} R_l(x)$$

$$\text{succ}(x^\omega) \stackrel{\text{def}}{=} \text{succ}(x)$$

$$\text{pred}(x^\omega) \stackrel{\text{def}}{=} \text{pred}(x)$$

4 Substructures

- in the example discussed (*ball*), we could also represent double l as a **substructure** if the model

$$\boxed{\phi_u(x) \stackrel{\text{def}}{=} l(x) \wedge l(\text{succ}(x))}$$

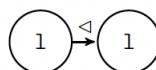
- in $\mathcal{M}_{ball}^{\triangleleft}$, $\phi_u(x)$ is true for $x = 3$
- substructures are formalized as Existentially Quantified Conjunctions (EQCs)
- with the domain $\mathcal{D} = \{1, 2, \dots, n\}$ and the set of relations $\mathcal{R} = \{R_1, R_2, \dots, R_m\}$, an EQC is the conjunction of expressions that specify which positions belong to which relations

$$\boxed{(\exists x_1 \dots, x_n) [\bigwedge_{i=1}^m R_i(x_a, \dots, x_b)]}$$

- example:

$$\boxed{\phi_u(x) \stackrel{\text{def}}{=} (\exists x, y) [l(x) \wedge l(\text{succ}(x)) \wedge x \triangleleft y]}$$

Visual representation of the EQC $\phi_u(x)$:



5 Formal Languages and Substructure Constraints

- we can use **formal languages** (set of well-formed strings) to represent SRs
- let \mathcal{K} be the logical formula, Σ the alphabet, \mathfrak{M} the model theory, w a word in Σ^* , and \mathcal{M}_w the model of w . Then \mathcal{K} defines a formal language ($L_{\mathcal{K}}$) which is a set of words in Σ^* whose word models satisfy. \mathcal{K}

$$L(\kappa) \stackrel{\text{def}}{=} \{w \in \Sigma^* \mid \mathcal{M}_w \models \kappa\}$$

- example:

If expression:

$$\psi_d \stackrel{\text{def}}{=} (\exists x)[d(x)]$$

and a **substructure constraint** \rightarrow conjunction of one or more literal. This could be though of as a well-formedness constraint:

$$\kappa_{ball} \stackrel{\text{def}}{=} \psi_{11} \wedge \neg\psi_d$$

Then the language $L(\kappa_{ball})$ consists of word models that contain two consecutive ll and no d :

$$L(\kappa_{ball}) = \{w \in \Sigma^* \mid \mathcal{M}_w^{\triangleleft} \models (\psi_{11} \wedge \neg\psi_d)\}$$

- if N is a logical formula that is the conjunction of all substructure constraints needed to define a particular pattern, $L(N)$ is the set of all words generated by the **substructure constraint grammar** N .

6 User-defined Formulas

- x is b: $b(x) \stackrel{\text{def}}{=} \text{voice}(x) \wedge \text{lab}(x) \wedge \text{stop}(x)$
- x is an obstruent: $\text{obs}(x) \stackrel{\text{def}}{=} \text{stop}(x) \vee \text{fric}(x)$
- x is a sonorant: $\text{son}(x) \stackrel{\text{def}}{=} \neg\text{obs}(x)$
- positions:

$$\begin{aligned} \text{init}(x) &\stackrel{\text{def}}{=} \text{pred}(x) = x \\ \text{fin}(x) &\stackrel{\text{def}}{=} \text{succ}(x) = x \\ \text{med}(x) &\stackrel{\text{def}}{=} \neg(\text{init}(x) \vee \text{fin}(x)) \end{aligned}$$