Perception-based Grammatical Inference for Adaptive Systems

Jie Fu (U. Pennsylvania)

Jeffrey Heinz (Delaware)

Adam Jardine (Delaware)

Herbert G. Tanner (Delaware)

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This paper

- 1. We introduce a learning paradigm called sensor-identification in the limit from positive data.
- 2. sensor is a perception module that obfuscates the learner's input.
- 3. Exact identification is eschewed for converging to a grammar which generates a language approximating the target language.
- 4. Successful approximation is understood as matching up to observation-equivalence.
- 5. Theoretical work exists which addresses other kinds of imperfect presentations, oracles, and the kinds of results obtainable with them [AL88, Ste95, FJ96, CJ01, THJ06].

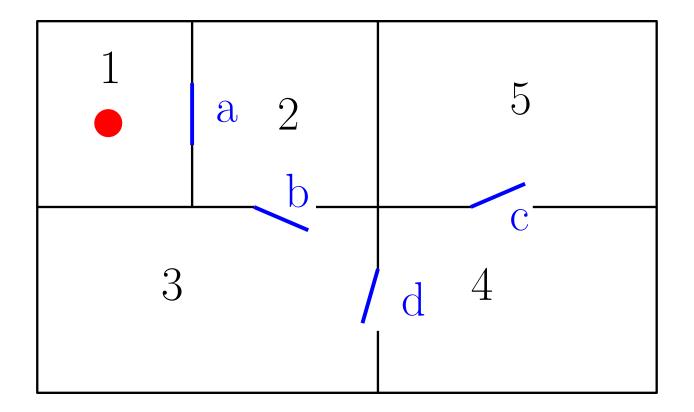
Motivation (part I)

- 1. A frontier in robotics is managing uncertainty.
- 2. Earlier work showed how to use grammatical inference to reduce the uncertainty in environments with potentially adversarial, but rule-governed behavior [CFK+12, FTH13, FTHC14].
- 3. The robot's capabilities, task, and environment were modeled as finite-state transition systems and product operations brought these elements together to form a *game*, allowing optimal control strategies to be computed (if they exist).
- 4. However, that work assumed *perfect* information about the environment.

Motivation (part II)

- 1. Recent results in game theory [AVW03, CDHR06] shows that optimal strategies can be found even for games with *imperfect* information (where players only have partial information about the state of the game).
- 2. The techniques in [CFK⁺12, FTH13, FTHC14] allow *imperfect* games to be constructed from *imperfect*—but consistent—models of the environment.
- 3. What is missing then is a way to identify such models from imperfect observations.
- 4. (POMDPs and MDPs address 1-player stochastic games, not 2-player games.)

Motivating Example



Basic Strategy

- 1. Convert learning solutions in the identification in the limit from positive data paradigm to solutions in the sensor-identification paradigm.
- 2. We focus on learnable regular classes of languages, which are well-studied [dlH10].

Sensor models

Sensor models have been proposed [CL08, LEPDG11, FDT14]. The definition below subsumes them all.

A sensor model is sensor = $\langle \Theta, \Sigma, \sim_{\theta} (\forall \theta \in \Theta), L_{\Theta} \rangle$ where

- Θ and Σ are finite, ordered sets of alphabets (the former being the sensor configurations).
- For all $\theta \in \Theta$, \sim_{θ} is an equivalence relation on Σ . If $\sigma_1 \sim_{\theta} \sigma_2$ then σ_1 is *indistinguishable* from σ_2 under sensor configuration θ . Let $[\sigma]_{\theta} = {\sigma' \in \Sigma \mid \sigma' \sim_{\theta} \sigma}$.
- $L_{\Theta} \subseteq \Theta^*$ is regular and represents the permissible sequences of sensor configurations.

We let $\hat{\Sigma}$ denote the powerset of Σ . So $[\sigma]_{\theta} \in \hat{\Sigma}$.

Observations (part I)

- 1. A bi-word is an element of $(\Theta \times \Sigma)^*$.
- 2. Let π_1 and π_2 be the left and right projections of $w \in (\Theta \times \Sigma)^*$.
- 3. obs: $(\Theta \times \Sigma)^* \to \hat{\Sigma}^*$ is defined inductively as follows.
 - The base case: $obs(\lambda) = {\lambda}$.
 - The inductive case:

$$\mathsf{obs}(w \cdot (\theta, \sigma)) = \mathsf{obs}(w) \cdot [\sigma]_{\theta}$$

4. Thus obs(u, v) is the *finite* set of sequences in Σ^* that are indistinguishable from v given the sequence u of sensor configurations.

Running Example (1)

Let
$$\Theta = \{\theta\}, \Sigma = \{0, 1, 2\}, \text{ and } [0]_{\theta} = \{0\} \text{ and } [1]_{\theta} = [2]_{\theta} = \{1, 2\}.$$

Consider the biword $w = (\theta, 0)(\theta, 1)(\theta, 1)(\theta, 0)(\theta, 2)(\theta, 2)$.

Then:

1.
$$\pi_1(w) = \theta\theta\theta\theta\theta\theta\theta$$
.

2.
$$\pi_2(w) = 011022$$
.

3.

$$\begin{aligned} \mathsf{obs}(w) &= & [0]_\theta \ [1]_\theta \ [0]_\theta \ [2]_\theta \ [2]_\theta \\ &= & \{0\}\{1,2\}\{1,2\}\{0\}\{1,2\}\{1,2\} \end{aligned}$$

Observations (part II)

Similarly, each $u \in \Theta^*$, a sensor model inductively induces an equivalence relation \sim_u over Σ^* .

- The base case: $\lambda \sim_{\lambda} \lambda$
- The inductive case: $(\forall \sigma_1, \sigma_2 \in \Sigma, v_1, v_2 \in \Sigma^*, \theta \in \Theta, u \in \Theta^*)$

$$[v_1 \sim_u v_2 \Rightarrow (v_1 \sigma_1 \sim_{u\theta} v_2 \sigma_2 \Leftrightarrow \sigma_1 \sim_{\theta} \sigma_2)]$$

Let $[v]_u = \{v' \in \Sigma^* \mid v \sim_u v'\}$, which denotes equivalent strings in Σ^* according to $u \in \Theta^*$.

Lemma 1. For all $w \in (\Theta \times \Sigma)^*$, $[\pi_2(w)]_{\pi_1(w)} = \mathsf{obs}(w)$ is a finite subset of Σ^* .

Running Example (2)

Consider biwords

$$w_1 = (\theta, 0)(\theta, 1)(\theta, 1)(\theta, 0)(\theta, 2)(\theta, 2)$$

$$w_2 = (\theta, 0)(\theta, 2)(\theta, 1)(\theta, 0)(\theta, 1)(\theta, 2)$$

Then

- 1. $obs(w_1) = obs(w_2)$
- 2. $w_1 \sim_{\theta\theta\theta\theta\theta\theta} w_2$

Facts and Observations

Facts on the Ground

Given L_{Θ} and L_{Σ} , the facts on the ground are

$$L_{\mathsf{system}} \stackrel{\text{def}}{=} \left\{ w \in (\Theta \times \Sigma)^* \mid \pi_1(w) \in L_\Theta \text{ and } \pi_2(w) \in L_\Sigma \right\}$$

The Observations on the Ground

In contrast, the observations on the ground are:

$$L_{\mathsf{sensor}} \stackrel{\text{def}}{=} \left\{ \hat{w} \in (\Theta \times \hat{\Sigma})^* \mid \exists w \in L_{\mathsf{system}} \text{ and} \right.$$
$$\pi_1(\hat{w}) = \pi_1(w) \text{ and } \pi_2(\hat{w}) = \mathsf{obs}(w) \right\}$$

Running Example (3)

Consider the languages

$$L_{\Theta} = \theta^*$$

$$L_{\Sigma} = \{ w \mid |w|_0, |w|_1, |w|_2 \text{ are each even} \}$$

Then

1.
$$w_1 = (\theta, 0)(\theta, 1)(\theta, 1)(\theta, 0)(\theta, 2)(\theta, 2)$$
 and $w_2 = (\theta, 0)(\theta, 2)(\theta, 1)(\theta, 0)(\theta, 1)(\theta, 2)$ belong to L_{system} .

2.
$$\left(\theta, \{0\}\right)\left(\theta, \{1, 2\}\right)\left(\theta, \{1, 2\}\right)\left(\theta, \{0\}\right)\left(\theta, \{1, 2\}\right)\left(\theta, \{1, 2\}\right)$$
 is an element of L_{sensor} .

Observation-equivalence of Languages

Definition 1 (Observation-equivalence). According to model sensor, languages $L, L' \subseteq \Sigma^*$ are observation-equivalent if

$$(\forall v \in L)(\exists v' \in L')(\forall u \in \{u \mid (u, v) \in L_{\mathsf{system}}\}) [v \sim_u v'] \qquad and$$
$$(\forall v' \in L')(\exists v \in L)(\forall u \in \{u \mid (u, v') \in L'_{\mathsf{system}}\}) [v \sim_u v']$$

Running Example (4)

Fix $L_{\Theta} = \theta^*$. Consider $L_t = \{ w \mid |w|_0, |w|_1, |w|_2 \text{ are each even} \}$ $L_h = \{ w \mid |w|_0, (|w|_1 + |w|_2) \text{ are both even} \}$

Then

1. L_t is observation-equivalent to L_h .

Illustration: Let $w_3 = (\theta, 1)(\theta, 1)(\theta, 1)(\theta, 2)(\theta, 2)(\theta, 2)$.

Then $\pi_2(w_3) = 111222 \in L_h$ but $\pi_2(w_3) \not\in L_t$. Nonetheless, $\mathsf{obs}(w_3) = \{1, 2\}\{1, 2\}\{1, 2\}\{1, 2\}\{1, 2\}\{1, 2\}\{1, 2\}\}$ and there exists w_4 such that $\pi_2(w_4) = 112211 \in L_t$ such that $\mathsf{obs}(w_4) = \mathsf{obs}(w_3)$.

Sensor-identification in the limit

We consider a sensor model sensor = $\langle \Theta, \Sigma, \sim_{\theta} (\forall \theta \in \Theta), L_{\Theta} \rangle$ and family of languages \mathcal{L} over Σ .

 \mathcal{L} is sensor-identifiable in the limit from positive data if there exists an algorithm \mathfrak{A} such that for all $L \in \mathcal{L}$, for any presentation ϕ of L_{sensor} , there exists $n \in \mathbb{N}$ such that for all $m \geq n$,

- $\mathfrak{A}(\phi[m]) = \mathfrak{A}(\phi[n]) = G$, and (convergence)
- L(G) is observation-equivalent to L. ("correctness")

Running Example (5)

If the target language is this one:

$$L_t = \{ w \mid |w|_0, |w|_1, |w|_2 \text{ are each even} \}$$

Then presentations draw elements from L_{sensor} :

not
$$(\theta, 0)(\theta, 0)(\theta, 1)(\theta, 1)(\theta, 2)(\theta, 2)$$
 but
$$(\theta, \{0\}) (\theta, \{0\}) (\theta, \{1, 2\}) (\theta, \{1, 2\}) (\theta, \{1, 2\}) (\theta, \{1, 2\}) (\theta, \{1, 2\})$$

not
$$(\theta, 1)(\theta, 0)(\theta, 2)(\theta, 0)(\theta, 1)(\theta, 2)$$
 but
$$(\theta, \{1, 2\}) (\theta, \{0\}) (\theta, \{1, 2\}) (\theta, \{0\}) (\theta, \{1, 2\}) (\theta, \{1, 2\}$$

. . .

Learning regular languages

For any L, let \sim_L be the Myhill-Nerode equivalence relation for L.

$$w \sim_L w' \Leftrightarrow \{v \in \Sigma^* \mid wv \in L\} = \{v \in \Sigma^* \mid w'v \in L\}$$

.

- 1. Given as input a finite sample $S \subset \Sigma^*$, a learning algorithm \mathfrak{A} determines an equivalence relation $\sim_{\mathfrak{A}}$ over Σ^* .
- 2. For any regular L, for any presentation ϕ of L, if $\mathfrak{A}(\phi)$ outputs $\sim_{\mathfrak{A}}$, which is of finite index and refines \sim_L then \mathfrak{A} identifies L in the limit from positive data.
- 3. If \mathfrak{A} does this for every $L \in \mathcal{L}$ then \mathfrak{A} identifies \mathcal{L} in the limit from positive data.

Useful Lemma

Lemma 2. If L_{Θ} and L are regular then $\sim_{L_{\mathsf{system}}}$ is of finite index and a right congruence. Furthermore,

$$w \sim_{\mathsf{system}} w' \Leftrightarrow \pi_1(w) \sim_{L_{\Theta}} \pi_1(w') \wedge \pi_2(w) \sim_L \pi_2(w')$$

Lifting congruences to $\hat{\Sigma}^*$

A right congruence \sim over Σ^* induces a relation \approx among elements of $\mathscr{P}(\Sigma^*)$:

$$X \approx Y \Leftrightarrow (\forall x \in X)(\exists y \in Y)(x \sim y) \land (\forall y \in Y)(\exists x \in X)[x \sim y]$$

Since elements of $\hat{\Sigma}^*$ can be understood as subsets of Σ^* , \approx_L is meaningful on $\hat{\Sigma}^*$.

Lemma 3. If \sim_{system} is of finite index and a right congruence then so is \sim_{sensor} . Furthermore,

$$w \sim_{\mathsf{sensor}} w' \Leftrightarrow \pi_1(w) \sim_{L_{\Theta}} \pi_1(w') \wedge \pi_2(w) \approx_L \pi_2(w')$$

- 1. By Lemmas 2 and 3, there is a DFA A accepting L_{sensor} . A defines a class of languages $\mathcal{L}_{\mathsf{sensor}}$ over Σ .
- 2. Each $L \in \mathcal{L}_{sensor}$ is obtained by replacing each label (which is an element of $\Theta \times \hat{\Sigma}$) of each transition in A with one element drawn from the label's right projection (thus the drawn element belongs to Σ).
- 3. These choices can be made consistently since Σ is ordered.

Lemma 4. Any $L' \in \mathcal{L}_{sensor}$ is observation-equivalent to L.

Main result

Theorem 1. Let \mathcal{L} be identifiable in the limit from positive data by a state-merging algorithm \mathfrak{A} and consider

sensor = $\langle \Theta, \Sigma, \sim_{\theta} (\forall \theta \in \Theta), L_{\Theta} \rangle$. There exists an algorithm \mathfrak{B} which Sensor-identifies \mathcal{L} in the limit from positive data.

Proof Sketch Algorithm \mathfrak{B} which takes as input a finite set $S \subset L_{\mathsf{sensor}}$ is defined from \mathfrak{A} which identifies \mathcal{L} , the equivalence relations $\theta \in \Theta$ on Σ , and L_{Θ} .

 \mathfrak{B} builds a PTA for S and merges prefixes according to $\sim_{\mathfrak{B}}$, defined as follows:

$$\hat{w} \sim_{\mathfrak{B}} \hat{w}' \Leftrightarrow \pi_1(\hat{w}) \sim_{L_{\Theta}} \pi_1(\hat{w}') \wedge \pi_2(\hat{w}) \approx_{\mathfrak{A}} \pi_2(\hat{w}').$$

(continued...)

Proof sketch (con't)

$$\hat{w} \sim_{\mathfrak{B}} \hat{w}' \Leftrightarrow \pi_1(\hat{w}) \sim_{L_{\Theta}} \pi_1(\hat{w}') \wedge \pi_2(\hat{w}) \approx_{\mathfrak{A}} \pi_2(\hat{w}').$$

- 1. Since L_{Θ} is regular, we assume it is given in terms of its minimal DFA and so $\sim_{L_{\Theta}}$ can be computed.
- 2. Also, $\approx_{\mathfrak{A}}$ can be computed since $\sim_{\mathfrak{A}}$ can be computed and every $\mathsf{obs}(w)$ ($w \in L_{\mathsf{system}}$) is a finite set.
- 3. In the limit, $\sim_{\mathfrak{B}}$ is of finite index because $\sim_{\mathfrak{A}}$ is of finite index.
- 4. Also in the limit, $\sim_{\mathfrak{B}}$ refines \sim_{sensor} because $\sim_{\mathfrak{A}}$ refines \sim_L in the limit and by definition of \approx .
- 5. Thus this acceptor recognizes the same language as L_{sensor} , and by Lemma 4, a language L' observation-equivalent to L can be obtained.
- 6. Convergence to L' is guaranteed by drawing least elements to find it.

Demonstration #1: Zero-reversible languages \mathcal{L}_{ZR}

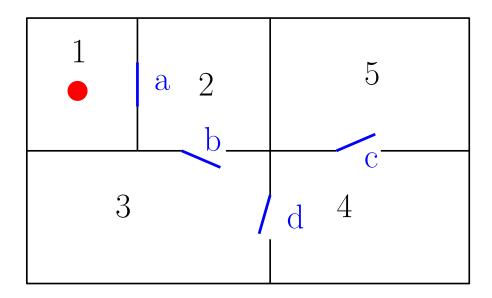
$$L_t = \{ w \mid |w|_0, |w|_1, |w|_2 \text{ are each even} \} \in \mathcal{L}_{ZR}$$

With a sufficient sample, \mathfrak{B} outputs a DFA recognizing this language.

$$L_h = \{ w \mid |w|_0 \text{ and } (|w|_1 + |w|_2) \text{ are both even} \}$$

As mentioned, this hypothesized language L_h is observation-equivalent to the target L_t .

Demonstration #2: Robot motion planning



- 1. The game is turn based. The robot can only move to an adjacent room if the adjoining door is open.
- 2. The dynamic, adversarial environment opens and closes doors according to a Strictly 2-Local language. For instance perhaps the same door cannot be closed on consecutive terms.
- 3. The robot can only see which doors are open/closed which adjoin to the room it is in.

Conclusion

- 1. Using the aforementioned strategy, an observation-equivalent language can be learned.
- 2. Techniques described in [CFK⁺12, FTH13, FTHC14] allow an imperfect game to be constructed.
- 3. Techniques from algorithmic game theory [AVW03, CDHR06] allow optimal strategies to be found.
- 4. Consequently, robots can deal with uncertainty better than before.

Thank you.

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A. D. Elandon D. Donamantha and I.A. De Claude Maltinal at an